

Appl. No. 09/918,831  
Amendment and/or Response  
Reply to Office action of 5 May 2005

Page 2 of 8

**Amendments to the Claims:**

A listing of the entire set of pending claims (including amendments to the claims, if any) is submitted herewith per 37 CFR 1.121. This listing of claims will replace all prior versions, and listings, of claims in the application.

**Listing of Claims:**

1. (Previously presented) A method of generating a linear transformation matrix A for use in a symmetric-key cipher, the method including:

generating a binary  $[n,k,d]$  error-correcting code, represented by a generator matrix  $G \in Z_2^{k \times n}$  in a standard form  $G = (I_k \parallel B)$ , with  $B \in Z_2^{k \times (n-k)}$ , where  $k < n < 2k$ , and d is the minimum distance of the binary error-correcting code;

extending matrix B with  $2k-n$  columns such that a resulting matrix C is non-singular, and

deriving matrix A from matrix C.

2. (Previously presented) A method as claimed in claim 1, wherein extending matrix B with  $2k-n$  columns includes:

in an iterative manner:

randomly generating  $2k-n$  columns, each with k binary elements;

forming a test matrix consisting of the  $n-k$  columns of B and the  $2k-n$  generated columns; and

checking whether the test matrix is non-singular,

until a non-singular test matrix has been found; and

using the found test matrix as matrix C.

Appl. No. 09/918,831  
Amendment and/or Response  
Reply to Office action of 5 May 2005

Page 3 of 8

3. (Previously presented) A method as claimed in claim 1, wherein deriving matrix A from matrix C includes:

determining two permutation matrices  $P_1, P_2 \in Z_2^{k \times k}$  such that all codewords in an  $[2k, k, d]$  error-correcting code, represented by the generator matrix  $(I \parallel P_1 C P_2)$ , have a predetermined multi-bit weight; and

using  $P_1 C P_2$  as matrix A.

4. (Original) A method as claimed in claim 3, wherein the cipher includes a round function with an S-box layer with S-boxes operating on m-bit sub-blocks, and the minimum predetermined multi-bit weight over all non-zero codewords equals a predetermined m-bit weight.

5. (Previously presented) A method as claimed in claim 3, wherein determining the two permutation matrices  $P_1$  and  $P_2$  includes iteratively generating the matrices in a random manner.

6. (Original) A method as claimed in claim 1, wherein the cipher includes a round function operating on 32-bit blocks and wherein the step of generating a  $[n, k, d]$  error-correcting code includes:

generating a binary extended Bose-Chaudhuri-Hocquenghem (XBCH) [64, 36, 12] code; and

shortening this code to a [60, 32, 12] shortened XBCH code by deleting four rows.

7. (Original) A computer program product, wherein the program product is operative to cause a processor to perform the method of claim 1.

**Appl. No. 09/918,831  
Amendment and/or Response  
Reply to Office action of 5 May 2005**

**Page 4 of 8**

8. (Previously presented) A system for cryptographically converting an input data block into an output data block; the data blocks comprising  $n$  data bits; the system including:

- an input for receiving the input data block;
- a storage for storing a linear transformation matrix  $A$ , generated according to the method of claim 1,
- a cryptographic processor performing a linear transformation on the input data block or a derivative of the input data block using the linear transformation matrix  $A$ ;
- and
- an output for outputting the processed input data block.

9-10 (Canceled)